



Written by :

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**Msc:Mathematics**

Notes:

Fuzzy sets theory

Followed book:

Mathematics of fuzzy sets and fuzzy logic

For Students:

For MSc and BS

Dedicated to:

All my MSc friends

Spacial thanks to:

Fazal Ullah Fazal (Msc mathematics )

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Fuzzy set :- Let  $X$  is non empty set ( $X \neq \phi$ ) then a fuzzy set is defined as follows:

$$A = \{ (x, \mu_A(x)) \mid x \in X \}$$

Where  $x$  is a particular element of  $X$  and  $\mu_A: X \rightarrow [0, 1]$  is called membership function and  $\mu_A(x)$  is called degree of membership of  $x$ .

Example :-  $X = \{1, 2, 3, 4, 5\}$

then fuzzy set of  $X$  as:

$$A = \{ (1, 0.9), (2, 0.5), (3, 0.4), (4, 0.6), (5, 0.2) \}$$

or also we written as

$$\mu_A(1) = 0.9, \quad \mu_A(2) = 0.5$$

$$\mu_A(3) = 0.4, \quad \mu_A(4) = 0.6, \quad \mu_A(5) = 0.2$$

Lotfi Zadeh :-

- \* He was born in 1921 in Azerbaijan (His family from Iran).
- \* He was mathematician, electric engineer and artificial intelligence teacher.
- \* professor of computer science at the university of California, Berkeley (US) (Amerika).
- \* He published fuzzy set theory in 1965.
- \* In 1973 he proposed his theory of fuzzy logic.
- \* Died in 6 september 2017.

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### Intersection of fuzzy sets:-

Let  $U_A$  and  $U_B$  be two fuzzy sets of  $X$ . Then intersection of  $U_A$  and  $U_B$  is denoted by  $U_A \cap U_B$  and defined as

$$(U_A \cap U_B)(x) = \min\{U_A(x), U_B(x)\}$$

Example :-

$$X = \{1, 2, 3\}$$

$$U_A = \{(1, 0.8) (2, 0.3) (3, 0.7)\}$$

$$U_B = \{(1, 0.2) (2, 0.7) (3, 0.3)\}$$

$$(U_A \cap U_B) = \{(1, \min(0.8, 0.2)), (2, \min(0.3, 0.7))$$
$$(3, \min(0.7, 0.3))\}$$

$$= \{(1, 0.2) (2, 0.3) (3, 0.3)\}$$

Example :-

$$X = \{x_1, x_2\}$$

$$U_A = \{(x_1, 0.3) (x_2, 0.8)\}$$

$$U_B = \{(x_1, 0.7) (x_2, 0.2)\}$$

$$(U_A \cap U_B) = \{(x_1, 0.3) (x_2, 0.2)\}$$

### Union of fuzzy sets:-

Let  $U_A$  and  $U_B$  be two fuzzy sets of  $X$ . Then Union of  $U_A$  and  $U_B$  is denoted by  $U_A \cup U_B$  and defined as

$$(U_A \cup U_B)(x) = \max\{U_A(x), U_B(x)\}$$

Example :-  $X = \{1, 2, 3\}$

$$U_A = \{(1, 0.8) (2, 0.3) (3, 0.7)\}$$

$$U_B = \{(1, 0.2) (2, 0.7) (3, 0.3)\}$$

$$(U_A \cup U_B)(x) = \{(1, \max(0.8, 0.2)) (2, \max(0.3, 0.7)) (3, \max(0.7, 0.3))\}$$

$$= \{(1, 0.8) (2, 0.7) (3, 0.7)\}$$

Example :-

$$X = \{x_1, x_2\}$$

$$U_A = \{(x_1, 0.3) (x_2, 0.8)\}$$

$$U_B = \{(x_1, 0.7) (x_2, 0.2)\}$$

$$(U_A \cup U_B)(x) = \{(x_1, 0.7), (x_2, 0.8)\}$$

Universal set :- A set whose degree of membership is 1 called universal set.

empty set :- A set whose degree of membership is 0 called empty set.

Compliment :- Compliment of fuzzy set is denoted by  $U_A^c(x)$  and defined as:

$$U_A^c(x) = 1 - U_A(x)$$

Example :-  $X = \{1, 2, 3\}$

$$A = \{(1, 0.6) (2, 0.3) (3, 0.7)\}$$

then

$$(A^c)(x) = \{(1, 1-0.6), (2, 1-0.3), (3, 1-0.7)\}$$

$$= \{(1, 0.4) (2, 0.7) (3, 0.3)\}$$

Example:-  $X = \{x_1, x_2, x_3\}$

$A = \{(x_1, 0.5) (x_2, 0.6) (x_3, 0.4)\}$

$A^c = \{(x_1, 0.5) (x_2, 0.4) (x_3, 0.6)\}$

Cross product (Cartesian product):-

Let  $U_1, U_2, U_3, \dots, U_n$  be  $n$ -fuzzy sets of  $X_1, X_2, X_3, \dots, X_n$  respectively,

$$U_1 \times U_2 \times U_3 \times \dots \times U_n (x_1, x_2, x_3, \dots, x_n) = \min \{U_1(x_1), U_2(x_2), U_3(x_3), \dots, U_n(x_n)\}$$

Example:-

$$X_1 = \{1, 2, 3\}$$

$$X_2 = \{4, 5, 6\}$$

$$A = \{(1, 0.2) (2, 0.3) (3, 0.5)\}$$

$$B(x) = \{(4, 0.8) (5, 0.6) (6, 0.3)\}$$

$$U_1 \times U_2 (x_1, x_2) = \min (A(x_1), B(x_2))$$

		4	5	6
A x B	1	0.2	0.2	0.2
	2	0.3	0.3	0.3
	3	0.5	0.5	0.3

$$= \{ [(1, 4), 0.2], [(1, 5), 0.2], [(1, 6), 0.2], [(2, 4), 0.3], [(2, 5), 0.3], [(2, 6), 0.3], [(3, 4), 0.5], [(3, 5), 0.5], [(3, 6), 0.3] \}$$

Example  $X_1 = \{1, 2, 3, 4\}$

$X_2 = \{5, 6, 7, 8\}$

$A(X_1) = \{(1, 0.3) (2, 0.5) (3, 0.7) (4, 0.9)\}$

$B(X_2) = \{(5, 0.2) (6, 0.4) (7, 0.6) (8, 0.8)\}$

$(A \times B)(X_1, X_2) = \min \{u_1(x_1), u_2(x_2)\}$

$A \times B = \{ [(1, 5), \min(0.3, 0.2)] [(1, 6), \min(0.3, 0.4)]$

$[(1, 7), \min(0.3, 0.6)] [(1, 8), \min(0.3, 0.8)]$

$[(2, 5), \min(0.5, 0.2)] [(2, 6), \min(0.5, 0.4)] [(2, 7), \min(0.5, 0.6)]$

$[(2, 8), \min(0.5, 0.8)] [(3, 5), \min(0.7, 0.2)] [(3, 6), \min(0.7, 0.4)]$

$[(3, 7), \min(0.7, 0.6)] [(3, 8), \min(0.7, 0.8)] [(4, 5), \min(0.9, 0.2)]$

$[(4, 6), \min(0.9, 0.4)] [(4, 7), \min(0.9, 0.6)] [(4, 8), \min(0.9, 0.8)] \}$

$= \{ [(1, 5), 0.2], [(1, 6), 0.3] [(1, 7), 0.3] [(1, 8), 0.3]$

$[(2, 5), 0.2] [(2, 6), 0.4] [(2, 7), 0.5] [(2, 8), 0.5]$

$[(3, 5), 0.2] [(3, 6), 0.4] [(3, 7), 0.6] [(3, 8), 0.7]$

$[(4, 5), 0.2] [(4, 6), 0.4] [(4, 7), 0.6] [(4, 8), 0.8] \}$

Level set ( $\alpha$ -cutset) :-

Let  $X$  be a non empty set and

$U_A$  be a fuzzy set of  $X$

then level set or  $\alpha$ -cut set is defined as

$A_\alpha = \{x \in X \mid U_A(x) \geq \alpha\}$  for  $\alpha \in (0, 1]$

example :-

$X = \{1, 2, 3\}$

$U_A = \{(1, 0.3) (2, 0.7) (3, 0.5)\}$

$$\alpha = 0.3$$

then  $A_\alpha = \{x \in X \mid \mu \geq \alpha\}$

$$A_{0.3} = \{1, 2, 3\}$$

if  $\alpha = 0.6$

then  $A_{0.6} = \{2\}$

if  $\alpha = 0.4$

then  $A_{0.4} = \{2, 3\}$

**Core set** :- Let  $X$  be non-empty and  $\mu_A$  be a fuzzy of  $X$  then core set is defined as.

$$A_1 = \{x \in X \mid \mu_A(x) \geq 1\}$$

Example :-

$$X = \{x_1, x_2, x_3\}$$

$$\mu_A = \{(x_1, 1) (x_2, 1) (x_3, 0.9)\}$$

$$A_1 = \{x_1, x_2\}$$

**Support set** :- Let  $X$  be non-empty set and  $\mu_A$  be a fuzzy of  $X$  then support set be defined as.

$$\text{Supp } A = \{x \in X \mid \mu_A(x) > 0\}$$

Example :-

$$X = \{x_1, x_2, x_3\}$$

$$\mu_A = \{(x_1, 0) (x_2, 0.3) (x_3, 0.1)\}$$

then

$$\text{Supp } A = \{x_2, x_3\}$$

**Subset of Fuzzy sets** :-

Let  $\mu_A$  and  $\mu_B$  be two



fuzzy sets of  $X$  then we said that  $U_A$  is subset of  $U_B$  if  $U_A(x) \leq U_B(x)$   
 $\forall x \in X$

Example :-

$$X = \{x_1, x_2\}$$

$$U_A = \{(x_1, 0.2) (x_2, 0.3)\}$$

$$U_B = \{(x_1, 0.4) (x_2, 0.5)\}$$

$$U_C = \{(x_1, 0.4) (x_2, 0.1)\}$$

Thus  $U_A \leq U_B$  because  $U_A(x) \leq U_B(x)$   
 $\forall x \in X$ .

And

$$U_A \not\leq U_C \text{ because } U_A(x_2) \neq U_C(x_2)$$

$$\text{Also } U_C \not\leq U_A \text{ because } U_C(x_1) \neq U_A(x_1).$$

## Properties of Fuzzy sets :-

Note we denote Union by  $\vee$  and intersection by  $\wedge$

(1) Associative property :-

$$U_A \wedge (U_B \wedge U_C) = (U_A \wedge U_B) \wedge U_C$$

LHS

$$(U_A \wedge (U_B \wedge U_C))(x) = \min \{U_A(x), (U_B \wedge U_C)(x)\}$$

$$= \min \{U_A(x), \min \{U_B(x), U_C(x)\}\}$$

$$= \min \{U_A(x), U_B(x), U_C(x)\}$$

RHS

$$((U_A \wedge U_B) \wedge U_C)(x) = \min \{(U_A \wedge U_B)(x), U_C(x)\}$$

$$= \min \{U_A(x), U_B(x), U_C(x)\}$$

$$\text{RHS} = \text{LHS}$$

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$$\text{Similarly } U_A \vee (U_B \wedge U_C) = (U_A \vee U_B) \wedge U_C$$

(2) Commutative property :-

$$U_A \wedge U_B = U_B \wedge U_A$$

$$U_A \vee U_B = U_B \vee U_A$$

(3) Identity property :-

$$A \wedge X = A$$

$$A \vee \phi = A$$

(4) Absorption by  $\phi$  and  $X$

$$A \wedge \phi = \phi$$

$$A \vee X = X$$

(5) Idempotence :-

$$A \vee A = A$$

$$A \wedge A = A$$

Proof :-

$$A \wedge A = A$$

$$(A \wedge A)_x = \min(A(x), A(x))$$

$$= A(x)$$

$$A \wedge A = A$$

similarly  
all (4) and (5)  
properties

(6) De Morgan Laws :-

$$\overline{A \wedge B} = \bar{A} \vee \bar{B}$$

LHS

$$(\overline{A \wedge B})_x = 1 - (A \wedge B)_x$$

$$= 1 - \min(A(x), B(x))$$

$$= \max\{1 - A(x), 1 - B(x)\}$$

$$= \max\{\bar{A}(x), \bar{B}(x)\}$$

$$= \bar{A}(x) \vee \bar{B}(x)$$

$$\overline{A \wedge B} = \bar{A} \vee \bar{B}$$

Similarly  $\overline{A \vee B} = \bar{A} \wedge \bar{B}$ .

(7) Distributive :-

$$A \wedge (B \vee C) = (A \wedge B) \vee (A \wedge C)$$

$$A \vee (B \wedge C) = (A \vee B) \wedge (A \vee C)$$

(8) Involution :-

$$\overline{\bar{A}} = A$$

$$\bar{A}(x) = 1 - A(x)$$

$$\bar{\bar{A}}(x) = 1 - \bar{A}(x) = 1 - (1 - A(x))$$

$$= A(x)$$

$$\bar{\bar{A}} = A$$

(9) Absorption :-

$$A \wedge (A \vee B) = A$$

$$(A \wedge (A \vee B))(x) = \min(A(x), (A \vee B)(x))$$

$$= \min(A(x), \max(A(x), B(x)))$$

$$= \min(A(x), B(x) \text{ or } A(x))$$

$$= A(x)$$

$$A \wedge (A \vee B) = A$$

Similarly  $A \vee (A \wedge B) = A$

Assignment :- If  $A$  is a non-classical fuzzy set  $A: X \rightarrow [0, 1]$  (i.e. then there exist  $x \in X$  with  $A(x) \notin \{0, 1\}$  then

$$A \wedge \bar{A} \neq \emptyset$$

$$A \vee \bar{A} \neq X$$

Proof :-

if  $x \in X$  is such that  $0 < A(x) < 1$

then  $0 < \bar{A}(x) < 1$  and then

$$0 < A \wedge \bar{A}(x) < 1 \quad \text{and} \quad 0 < A \vee \bar{A}(x) < 1$$

For example :-

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$$X = \{1, 2\}$$

$$A = \{(1, 0.4) (2, 0.3)\} \quad \text{and}$$

$$\bar{A} = \{(1, 0.6) (2, 0.7)\}$$

then

$$A \wedge \bar{A} = \{(1, \min(0.4, 0.6)) (2, \min(0.3, 0.7))\}$$

$$A \wedge \bar{A} = \{(1, 0.4) (2, 0.3)\} \neq \phi$$

$$\boxed{A \wedge \bar{A} \neq \phi}$$

Also

$$A \vee \bar{A} = \{(1, \max(0.4, 0.6)) (2, \max(0.3, 0.7))\}$$

$$A \vee \bar{A} = \{(1, 0.6) (2, 0.7)\} \neq X$$

$$\Rightarrow \boxed{A \wedge \bar{A} \neq X}$$

**Negation :-** A function  $N: [0, 1] \rightarrow [0, 1]$  is called Negation if

\*  $N(0) = 1$  and  $N(1) = 0$  and  $N$  is non-increasing ( $x < y \Rightarrow N(x) > N(y)$ ).

**Strict negation :-**

A Negation is called strict negation if it is strictly decreasing  $x < y \Rightarrow N(x) > N(y)$

**Strong negation :-**

A strict negation is said to be strong negation if it is also involutive i.e.  $N(N(x)) = x$ .

**Example :-**

$$N(x) = 1 - x$$

\*  $N(0) = 1$ ,  $N(1) = 0$ .

and  $x < y$

$$-x > -y \Rightarrow 1 - x > 1 - y$$

$$\Rightarrow N(x) > N(y)$$

It is negation.

\*  $x < y \Rightarrow -x > -y$ .

$$\Rightarrow 1 - x > 1 - y \Rightarrow N(x) > N(y)$$

It is strict negation.

\*  $N(x) = 1 - x$

$$N(N(x)) = 1 - N(x)$$

$$= 1 - (1 - x)$$

$$= 1 - 1 + x$$

$$N(N(x)) = x$$

It is strong negation.

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Example :-  $N_\lambda(x) = \frac{1-x}{1+\lambda x}$ ,  $\lambda > -1$

\*  $N_\lambda(0) = 1$ ,  $N_\lambda(1) = 0$  and.

$$x \leq y$$

$$-x \geq -y \Rightarrow 1-x \geq 1-y$$

$$\Rightarrow \frac{1-x}{1+\lambda x} \geq \frac{1-y}{1+\lambda y}$$

$$\Rightarrow N_\lambda(x) \geq N_\lambda(y)$$

It is Negation.

\*  $x < y$

$$\frac{1-x}{1+\lambda x} > \frac{1-y}{1+\lambda y} \Rightarrow N_\lambda(x) > N_\lambda(y)$$

It is strict negation.

\*

$$N_\lambda(x) = \frac{1-x}{1+\lambda x}$$

$$N(N_\lambda(x)) = \frac{1 - N_\lambda(x)}{1 + \lambda N_\lambda(x)}$$

$$= \frac{1 - \frac{1-x}{1+\lambda x}}{1 + \lambda \left( \frac{1-x}{1+\lambda x} \right)} = \frac{1+\lambda x - (1-x)}{1+\lambda x + \lambda(1-x)}$$

$$= \frac{1+\lambda x - 1 + x}{1+\lambda x + \lambda - \lambda x} = \frac{\lambda x + x}{1+\lambda} = \frac{(\lambda+1)x}{1+\lambda}$$

$$N(N_\lambda(x)) = x$$

It is strong Negation

$\Rightarrow N_\lambda(x) = \frac{1-x}{1+\lambda x}$  strong negation called  $\lambda$ -complement.

## Triangular Norms and Conorms :-

### T-Norm (t-norm) :-

t-norm is function  $T: [0, 1] \times [0, 1] = [0, 1]^2 \rightarrow [0, 1]$

that satisfies the following properties:

$$T_1: T(x, 1) = x \text{ (identity)}$$

$$T_2: T(x, y) = T(y, x) \text{ (commutative)}$$

$$T_3: T(x, T(y, z)) = T(T(x, y), z) \text{ (associativity)}$$

$$T_4: \text{if } x \leq u \text{ and } y \leq v \text{ then} \\ T(x, y) \leq T(u, v)$$

### S-Norm (s-conorm) :-

s-norm is a function  $S: [0, 1]^2 \rightarrow [0, 1]$

that satisfies following properties:

$$S_1: S(x, 0) = x \text{ (identity)}$$

$$S_2: S(x, y) = S(y, x) \text{ (commutative)}$$

$$S_3: S(x, S(y, z)) = S(S(x, y), z)$$

associative

$$S_4: \text{if } x \leq u, y \leq v \text{ then} \\ S(x, y) \leq S(u, v)$$

Example :-  $T(x, y) = \min\{x, y\}$   $x, y \in [0, 1]$

$$T_1: T(x, 1) = \min\{x, 1\} = x \text{ satisfied.}$$

$$S_1: T(x, 0) = \min\{x, 0\} = 0 \neq x \text{ thus}$$

s-norm condition not

satisfied.

$$T_2: T(x, y) = \min\{x, y\} = \min\{y, x\} = T(y, x)$$

satisfied.

$$\begin{aligned}
 T_3: T\{u, T(y, z)\} &= \\
 &= \min\{u, \min(y, z)\} \\
 &= \min\{u, y, z\} \\
 &= \min\{\min(u, y), z\} \\
 &= \min\{T(u, y), z\} \\
 &= T\{T(u, y), z\} \text{ satisfy}
 \end{aligned}$$

$$\begin{aligned}
 T_4: u \leq u \text{ and } y \leq v \\
 \text{then } \min\{u, y\} \leq \min\{u, v\} \\
 \therefore T\{u, y\} \leq T\{u, v\}
 \end{aligned}$$

Thus

$$T(u, y) = \min\{u, y\} \text{ is t-norm.}$$

Example :-

$$T(u, y) = u \cdot y \text{ is t-norm?}$$

$$T_1: T\{u, 1\} = u \text{ satisfy}$$

$$T_2: T(u, y) = u \cdot y = y \cdot u = T(y, u) \text{ satisfy}$$

$$\begin{aligned}
 T_3: T(u, T(y, z)) &= T(u, y \cdot z) = u \cdot y \cdot z \\
 &= (u \cdot y) \cdot z = T(u \cdot y, z) \\
 &= T(T(u, y), z) \text{ satisfy}
 \end{aligned}$$

$$T_4: u \leq u \text{ and } y \leq v$$

$$\Rightarrow u \cdot y \leq u \cdot v$$

$$\Rightarrow T(u, y) \leq T(u, v)$$

$$T(u, y) = u \cdot y \text{ is t-norm proved.}$$

Example :-

$$S(u, y) = u + y - u \cdot y$$



$$S_1: S(x, 0) = x \quad \text{satisfy}$$

$$S_2: S(x, y) = x + y - xy$$

$$= y + x - yx$$

$$= S(y, x) \quad \text{satisfy}$$

$$S_3: S(x, S(y, z)) = S(x, y + z - yz)$$

$$\text{LHS} \rightarrow = x + y + z - yz - x(y + z - yz)$$

$$= x + y + z - yz - xy - xz + xyz$$

$$\text{RHS } S(S(x, y), z)$$

$$= S(x + y - xy, z)$$

$$= x + y - xy + z - (x + y - xy)z$$

$$= x + y - xy + z - xz - yz + xyz$$

$$= x + y + z - yz - xy - xz + xyz$$

$$\Rightarrow \text{LHS} = \text{RHS} \quad \text{satisfy}$$

$$S_4: x \leq u \quad \text{and} \quad y \leq v \quad \therefore x \in [0, 1]$$

$$(x + y - xy) \leq (u + v - uv)$$

$$S(x, y) \leq S(u, v)$$

example

$$0.2 < 0.3, \quad 0.4 < 0.5$$

$$0.2 + 0.4 - (0.2)(0.4) \leq 0.3 + 0.5 - (0.3)(0.5)$$

$$\boxed{0.52 < 0.65}$$

**Definition:-**

Let  $F(X)$  be collections of all fuzzy sets of  $X$ , and  $T, S: [0, 1]^2 \rightarrow [0, 1]$  at  $t$ -norm and  $t$ -conorm. then operation for any  $U_A, U_B$  can be defined as following;

$$(U_A T U_B)(x) = U_A(x) T U_B(x)$$

$$(U_A S U_B)(x) = U_A(x) S U_B(x)$$

**Proposition:- (2.11)**

Given any  $t$ -norm  $T$  and  $t$ -conorm  $S$  we have  $T(x, 0) = 0$  and  $S(x, 1) = 1 \forall x \in [0, 1]$

**Proof:-** From  $T_1$  we have

$$T(0, 1) = 0$$

$$T(0, x) \leq T(0, 1) = 0 \quad \therefore \text{From } T_4 \text{ as } 0 \leq 0, x \leq 1$$

$$\Rightarrow T(0, x) \leq 0$$

$$\Rightarrow T(0, x) = 0$$

$$\boxed{T(x, 0) = 0} \quad \therefore \text{From } T_2, T(x, y) = T(y, x)$$

Now From  $S_1$  we have

$$S(1, 0) = 1$$

$$1 = S(1, 0) \leq S(1, x) \quad \therefore \text{From } S_4 \text{ as } 1 \leq 1, 0 \leq x$$

$$\Rightarrow 1 \leq S(1, x)$$

$$1 = S(1, x)$$

$$\Rightarrow S(1, x) = 1$$

$$\boxed{S(x, 1) = 1} \quad \therefore \text{From } S_2, S(x, y) = S(y, x)$$

**Proposition:- (2.12)**

Let  $T$  and  $S$  be  $t$ -norm and  $t$ -conorm then,

- (i)  $T(x, y) \leq x \wedge y$  and (ii)  $S(x, y) \geq x \vee y$
- for any  $x, y \in [0, 1]$

**Proof:-** since from  $T_4$  we have

$$T(x, 1) = x$$

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$$T(x, y) \leq T(x, 1) = x \text{ --- ①} \because \text{From } T_4 \text{ as } x \leq x, y \leq 1 \\ \Rightarrow T(x, y) \leq T(x, 1)$$

Also From  $T_2$  we have

$$T(x, y) = T(y, x) \leq T(y, 1) = y \text{ --- ②}$$

$\therefore$  From  $T_4$ ,  $y \leq y, x \leq 1$

$$T(y, x) \leq T(y, 1)$$

then From  $T_1$ .

$$T(y, 1) = y$$

From ① and ②

$$T(x, y) \wedge T(x, y) \leq x \wedge y$$

$$\Rightarrow \boxed{T(x, y) \leq x \wedge y}$$

(ii)

$$x = S(x, 0) \leq S(x, y) \text{ --- ①} \because \text{From } S_4 \text{ and } S_1$$

Also

$$y = S(y, 0) \leq S(y, x) = S(x, y) \text{ --- ②}$$

$\therefore$  From  $S_2, S_4$  and  $S_1$

From ① and ②

$$S(x, y) \vee S(x, y) \geq x \vee y$$

$$\Rightarrow \boxed{S(x, y) \geq x \vee y}$$

De-Morgan Triplet :-

A triplet  $(S, T, N)$  is called De-Morgan triplet if  $T$  is t-norm and  $S$  is s-norm and  $N$  is a strong

Negation. Also Full fill De-Morgan law;

$$S(x, y) = N(T(N(x), N(y)))$$

2.15 (1)  
Example :-

ph 13

$$T(x, y) = x \wedge y$$

$$S(x, y) = x \vee y$$

$$N(x) = 1 - x$$

Form a de-Morgan triplet?

(1)  $T(x, y) = x \wedge y$  it is a t-norm  
see in ph 12 (example)

(2)  $S(x, y) = x \vee y$

(i)  $S(x, 0) = \max(x, 0) = \max(x, 0) = x$  satisfy.

(ii)  $S(x, y) = \max(x, y) = \max(y, x) = S(y, x)$

(iii)  $S(x, S(y, z)) = \max(x, y \vee z)$   
 $= \max(x, \max(y, z))$   
 $= \max(x, y, z)$

$S(S(x, y), z) = \max(x \vee y, z)$   
 $= \max(\max(x, y), z)$   
 $= \max(x, y, z)$

LHS = RHS

(iv) if  $x \leq u, y \leq v$

then  $\max(x, y) \leq \max(u, v)$

$S(x, y) \leq S(u, v)$

Thus it is s-norm.

(3)  $N(x) = 1 - x$  it is a strong  
negation see ph 10, example

(4)  $S(x, y) = N(T(N(x), N(y)))$   
 $= N[N(x) \wedge N(y)] \quad \because T\text{-norm}$   
 $= 1 - [N(x) \wedge N(y)] \quad \because \text{def Negation}$   
 $= (1 - N(x)) \vee (1 - N(y))$   
 $= (1 - (1 - x)) \vee (1 - (1 - y))$   
 $= x \vee y = S(x, y) \quad \because \text{def S-Norm}$

It form a de-Morgan triplet.

ph 19  
2.16

example (2) :-  $T(x, y) = x \cdot y$

$$S(x, y) = x + y - xy$$

$$N(x) = 1 - x \quad \text{form a}$$

Solution :-

de-morgan triplet

①  $T(x, y) = x \cdot y$

It is t-Norm (see ph 13)

②  $S(x, y) = x + y - xy$

It is S-Norm (see ph 13)

③  $N(x) = 1 - x$

It is strong negation

(see ph 10)

if given example satisfy de-morgan law then it form a de-morgan triplet.

$$S(x, y) = N[T(N(x), N(y))]$$

$$= N[N(x) \cdot N(y)] \quad \because \text{def } T\text{-Norm}$$

$$= 1 - (N(x))(N(y)) \quad \because \text{because of negation.}$$

$$= 1 - [(1-x)(1-y)]$$

$$= 1 - [1 - y - x + xy]$$

$$= 1 - 1 + y + x - xy$$

$$= x + y - xy$$

$$S(x, y) = S(x, y)$$

Thus it form de-morgan triplet.

example (3) (2.17) :-

$$T(x, y) = (x + y - 1) \vee 0$$

$$S(x, y) = (x + y) \wedge 1$$

$$N(x) = 1 - x$$

Form a demorgan triplet.

$$T(x, y) = \max[x + y - 1, 0] \quad \text{ph 20}$$

$$\text{(i)} T(x, 1) = \max[x + 1 - 1, 0] = \max[x, 0] \\ = x$$

$$\text{(ii)} T(x, y) = \max[x + y - 1, 0] = \max[0, x + y - 1] \\ = T(y, x)$$

$$\text{(iii)} T(x, T(y, z)) = \dots T(x, \max[y + z - 1, 0]) \\ = (x + (y + z - 1) - 1) \vee 0 \\ = (x + y + z - 2) \vee 0$$

$$T(T(x, y), z) = T((x + y - 1) \vee 0, z) \\ = (x + y - 1 + z - 1) \vee 0 \\ = (x + y + z - 2) \vee 0 \quad (\text{satisfies})$$

$$\text{(iv)} \quad x \leq u, \quad y \leq v$$

$$(x + y - 1) \vee 0 \leq (u + v - 1) \vee 0,$$

$$T(x, y) \leq T(u, v)$$

$$T(x, y) = (x + y - 1) \vee 0 \quad \text{is } t\text{-norm.}$$

$$\text{(i)} S(x, 0) = \dots [(x + y) \wedge 1] = \dots [x \wedge 1] \\ = x$$

$$\text{(ii)} S(x, y) = (x + y) \wedge 1 = (y + x) \wedge 1 = S(y, x)$$

$$\text{(iii)} S(x, S(y, z)) = \dots \\ = S(x, (y + z) \wedge 1) = (x + y + z) \wedge 1$$

$$S(S(x, y), z) = S((x + y) \wedge 1, z) \\ = (x + y + z) \wedge 1$$

$$\text{(iv)} \quad x \leq u, \quad y \leq v$$

$$(x + y) \wedge 1 \leq (u + v) \wedge 1$$

$$S(x, y) \leq S(u, v)$$

$$S(x, y) = (x + y) \wedge 1 \quad \text{is } S\text{-norm.}$$

ph 21

$N(x) = 1 - x$  is a strong negation

$$(4) S(x, y) = N[T(N(x), N(y))]$$

$$= N[(N(x) + N(y) - 1) \vee 0]$$

$$= N[(1-x + 1-y - 1) \vee 0]$$

$$= N[(1-x-y) \vee 0]$$

$$= 1 - [(1-x-y) \vee 0]$$

$$= 1 - \max[1-x-y, 0]$$

$$= \min[1 - (1-x-y), 1-0]$$

$$= \min[x+y, 1] = (x+y) \wedge 1 = S(x, y)$$

It is form a de-Morgan triplet.

Example (2.18):-

$$T(x, y) = \begin{cases} x \wedge y & \text{if } x+y > 1 \\ 0 & \text{if } x+y \leq 1 \end{cases}$$

$$(1) T(x, 1) = \begin{cases} x \wedge 1 & \text{if } x+1 > 1 \\ 0 & \text{if } x+1 \leq 1 \end{cases}$$

$$\text{Since } x+1 > 1$$

$$= \min(x, 1)$$

$$= x$$

satisfy.

$$(2) T(x, y) = \begin{cases} x \wedge y & \text{if } x+y > 1 \\ 0 & \text{if } x+y \leq 1 \end{cases}$$

$$= \begin{cases} y \wedge x & \text{if } y+x > 1 \\ 0 & \text{if } y+x \leq 1 \end{cases}$$

$$= T(y, x)$$

$$(3) T(x, T(y, z)) = \begin{cases} T(x, y \wedge z) & \text{if } y+z > 1 \\ 0 & \text{otherwise} \end{cases}$$

$$= \begin{cases} x \wedge y \wedge z & \text{if } x+y+z > 1 \\ 0 & \text{otherwise} \end{cases}$$

$$= \begin{cases} x \wedge y \wedge z & \text{if } x+y+z > 1 \\ 0 & \text{otherwise} \end{cases}$$

Also

$$T(T(x, y), z) = \begin{cases} T(x \wedge y, z) & \text{if } x+y > 1 \\ 0 & \text{otherwise} \end{cases}$$

$$= \begin{cases} x \wedge y \wedge z & \text{if } x \wedge y + z > 1 \\ 0 & \text{otherwise} \end{cases}$$

$$= \begin{cases} x \wedge y \wedge z & \text{if } x+y+z > 1 \\ 0 & \text{otherwise} \end{cases}$$

$$\Rightarrow T(x, T(y, z)) = T(T(x, y), z)$$

(4) if  $x \leq u, y \leq v$ .

Case (I)  $x+y > 1, u+v > 1$

$$x \wedge y \leq u \wedge v$$

$$T(x, y) \leq T(u, v)$$

Case (II)  $x+y \leq 1, u+v > 1$

$$0 \leq u \wedge y$$

$$T(x, y) \leq T(u, v)$$

Case (III)  $x+y \leq 1, u+v \leq 1$

$$0 \leq 0$$

$$T(x, y) \leq T(u, v)$$

$\Rightarrow T$  is t-norm.



ph 22.

$$S(x, y) = \begin{cases} xy & \text{if } x+y < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$(i) \quad S(x, 0) = \begin{cases} x \cdot 0 & \text{if } x+0 < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$\boxed{S(x, 0) = x}$$

$$(ii) \quad S(x, y) = \begin{cases} xy & \text{if } x+y < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$= \begin{cases} yx & \text{if } y+x < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$S(x, y) = S(y, x)$$

$$(iii) \quad S(x, S(y, z)) = \begin{cases} S(xy, z) & \text{if } y+z < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$= \begin{cases} xyvz & \text{if } x+yvz < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$= \begin{cases} xyvz & \text{if } x+y+z < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$S(S(x, y), z) = \begin{cases} S(xy, z) & \text{if } x+y < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$= \begin{cases} xyvz & \text{if } x+y+z < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$\Rightarrow S(x, S(y, z)) = S(S(x, y), z)$$

$$(iv) \quad x < u \quad y < v$$

$$x+y < 1 \quad y+u+v < 1$$

$$xy < uv \quad \Rightarrow S(x, y) < S(u, v)$$

Case II

$$x+y \leq 1 \text{ and } u+v \geq 1 \Rightarrow xy \leq 1$$

$$\max(u, v) \leq 1 \Rightarrow S(x, y) \leq S(u, v)$$

$$S(x, y) \leq S(u, v)$$

Case III  $x+y \geq 1$  and  $u+v \geq 1$ 

$$1 \leq 1$$

$$S(x, y) = S(u, v)$$

$\Rightarrow S(x, y)$  is t-conorm.

Now For de-morgan triplet

$$S(x, y) = N[T(N(x), N(y))]$$

$$= \begin{cases} N(N(x) \wedge N(y)) & \text{if } N(x) + N(y) > 1 \\ N(0) & \text{otherwise} \end{cases}$$

$$= \begin{cases} 1 - \min(N(x), N(y)) & \text{if } N(x) + N(y) > 1 \\ 1 - (0) & \text{otherwise} \end{cases}$$

$$= \begin{cases} \max(1 - N(x), 1 - N(y)) & \text{if } (1 - N(x)) + (1 - N(y)) < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$= \begin{cases} \max(1 - (1 - x), 1 - (1 - y)) & \text{if } (1 - (1 - x)) + (1 - (1 - y)) < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$= \begin{cases} \max(x, y) & \text{if } x + y < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$= \begin{cases} xy & \text{if } x + y < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$S(x, y) = S(x, y)$$

Thus triplet  $(S, T, N)$  Form a demorgan triplet

Fuzzy implication: Let  $I: [0,1]^2 \rightarrow [0,1]$  be a function if the following conditions are fulfilled:

(i) if  $x \leq y$ , then  $I(x,z) \geq I(y,z)$  i.e.  $I$  is decreasing in its first variable.

(ii) if  $y \leq z$  then  $I(x,y) \leq I(x,z)$  i.e.  $I$  is increasing in its second variable.

(iii)  $I(1,0) = 0$ ,  $I(0,0) = I(1,1) = 1$

then  $I$  is called fuzzy implication.

Example: (2.35)

$$I_3(x,y) = \begin{cases} 1 & \text{if } x \leq y \\ y/x & \text{if } x > y \end{cases}$$

are fuzzy implication.

proof:-

(i) if  $x \leq y$ , then  $I(x,z) \geq I(y,z)$ , i.e.  $I$  is decreasing in its first variable.

$\Rightarrow$  if  $x \leq y$ , there are three cases.

(a)  $x \leq y \leq z$

$$\Rightarrow x \leq z \Rightarrow I(x,z) = 1$$

$$\Rightarrow y \leq z \Rightarrow I(y,z) = 1$$

$$\text{Thus } \boxed{I(y,z) = I(x,z)}$$

(b)  $x \leq z < y$

$$\Rightarrow x \leq z \Rightarrow I(x,z) = 1$$

$$\Rightarrow y > z \Rightarrow I(y,z) = z/y$$

$$\text{Thus } \boxed{I(x,z) > I(y,z)}$$

(c)  $z < x < y$

$$\Rightarrow x > z \Rightarrow z/x \quad \text{since } \begin{matrix} x \leq y \\ x \geq z \end{matrix}$$

$$\Rightarrow y > z \Rightarrow z/y$$

$$\frac{z}{x} \geq \frac{z}{y}$$

$$\boxed{I(x,z) \geq I(y,z)}$$

Thus in all three cases  $I$  is decreasing in its 1st variable.

(ii) if  $y \leq z$  then  $I(n, y) \leq I(n, z)$   
i.e.  $I$  is increasing in 2nd variable  
if  $y \leq z$  then there are three cases

(a)  $x \leq y \leq z$

$$\Rightarrow x \leq y \Rightarrow I(n, y) = 1$$

$$\Rightarrow x \leq z \Rightarrow I(n, z) = 1$$

$$\text{Thus } \boxed{I(n, y) = I(n, z)}$$

(b)  $y < x \leq z$

$$\Rightarrow x \leq z \Rightarrow I(n, z) = 1$$

$$\Rightarrow y < x \Rightarrow I(n, y) = y/x$$

$$\text{Thus } \boxed{I(n, y) < I(n, z)}$$

(c)  $y < x < z$

$$\Rightarrow x > y \Rightarrow I(n, y) = y/x$$

$$\Rightarrow x > z \Rightarrow I(n, z) = z/x$$

Since  $y \leq z$

$$\frac{y}{x} \leq \frac{z}{x}$$

$$\Rightarrow \boxed{I(n, y) \leq I(n, z)}$$

Thus all in three cases  $I$  is increasing in 2nd variable.

(iii)  $I(1, 0) = 0$

$$\text{since } x > y \Rightarrow I(1, 0) = \frac{y}{x} = \frac{0}{1} = 0$$

$$\boxed{I(1, 0) = 0}$$

$$I(0, 0) = I(1, 1)$$

since  $x = y$  in both cases then

$$I(0,0) = I(1,1) = 1$$

Thus from (i), (ii) and (iii)  $I$  are fuzzy implication.

Example :-  $I(x,y) = \max\{1-x, y\}$  is fuzzy implication.

Proof :-

(i) if  $x \leq y$  then  $I(x,z) \geq I(y,z)$ .

$$x \leq y \Rightarrow I(x,z) \geq I(y,z)$$

$$\max\{1-x, z\} \geq \max\{1-y, z\}$$

example.

$$x = 0.3, y = 0.4, z = 0.5$$

$$\max\{1-x, z\} \geq \max\{1-y, z\}$$

$$\max\{0.7, 0.5\} \geq \max\{0.6, 0.5\}$$

$$\boxed{0.7 \geq 0.6}$$

$\Rightarrow$  since  $x \leq y$ .

$$-x \geq -y$$

$$1-x \geq 1-y$$

$$\max(1-x, z) \geq \max(1-y, z)$$

$$\boxed{I(x,z) \geq I(y,z)}$$

(ii)  $y \leq z$ .

$$\Rightarrow \max(1-x, y) \leq \max\{1-x, z\}$$

$$\boxed{I(x,y) \leq I(x,z)}$$

example  $0.4 \leq 0.6$

$$\max\{0.3, 0.4\} \leq \max\{0.3, 0.6\}$$

$$\boxed{0.4 \leq 0.6}$$

(iii)  $I(1,0) = \max\{1-1, 0\}$

$$= \max\{0, 0\} = 0$$

$$\Rightarrow \boxed{I(1,0) = 0}$$

$$I(0,0) = \max\{1-0, 0\} = \max\{1, 0\} = 1$$

$$\boxed{I(0,0) = 1}$$

$$I(1,1) = \max\{1-1, 1\} = \max\{0, 1\}$$

$$\Rightarrow \boxed{I(1,1) = 1}$$

Thus from (i) (ii) and (iii)  $I(x,y) = (1-x)v_y$  is fuzzy implication.

Example:  $I(x,y) = \min\{1-x+y, 1\}$

is fuzzy implication.

proof :-

(i) if  $x \leq y$

$$x \leq y$$

$$= -x \geq -y$$

$$= (1-x+z) \geq (1-y+z)$$

$$\min\{1-x+z, 1\} \geq \min\{1-y+z, 1\}$$

$$\boxed{I(x,z) \geq I(y,z)}$$

(ii) if

$$y \leq z$$

$$= -x+y \leq -x+z$$

$$= 1-x+y \leq 1-x+z$$

$$\min\{1-x+y, 1\} \leq \min\{1-x+z, 1\}$$

$$\boxed{I(x,y) \leq I(x,z)}$$

(iii)  $I(1,0) = \min\{1-1+0, 1\} = \min\{0, 1\}$

$$\Rightarrow \boxed{I(1,0) = 0}$$

$$I(1,1) = \min\{1-1+1, 1\} = \min\{1, 1\}$$

$$\boxed{I(1,1) = 1}$$

$$I(0,0) = \min\{1-0+0, 1\} = \min\{1, 1\}$$

$$\Rightarrow \boxed{I(0,0) = 1}$$

Thus from (i) (ii) and (iii)  $I$  is fuzzy implication.

Theorem 2.36:- If  $I$  is a fuzzy implication and if  $N$  is a negation then  $I'(x, y) = I(N(y), N(x))$  is a fuzzy implication.

Proof:-

(i) if  $x \leq y \Rightarrow N(y) \leq N(x)$   $\therefore$  by definition of negation.

Since  $I$  is implication then by using 2nd condition

if  $y \leq z$  then  $I(x, y) \leq I(x, z)$

since  $N(y) \leq N(x)$

then  $I(N(x), N(y)) \leq I(N(z), N(x))$

$I'(y, z) \leq I'(x, z)$

$\Rightarrow$  if  $x \leq y$  then  $I'(y, z) \leq I'(x, z)$

(ii) if  $y \leq z$

$= N(z) \leq N(y)$   $\therefore$  by definition of negation.

Since  $I$  is fuzzy implication then by using 1st condition.

if  $x \leq y \Rightarrow I(x, z) \geq I(y, z)$

since

$N(z) \leq N(y)$  or  $N(y) \geq N(z)$

$I(N(y), N(x)) \leq I(N(z), N(x))$

$I'(x, y) \leq I'(x, z)$

(iii)

$I'(1, 0) = I(N(0), N(1)) = I(1, 0) = 0$

$I'(1, 1) = I(N(1), N(1)) = I(0, 0) = 1$

$I'(0, 0) = I(N(0), N(0)) = I(1, 1) = 1$

Thus  $I'(x, y) = I(N(y), N(x))$  is a fuzzy implication.

$$(i) \quad x \leq y \Rightarrow N(x) \geq N(y)$$

$$S(N(x), z) \geq S(N(y), z)$$

$$I(x, z) \geq I(y, z)$$

ph 32  $I$  is decreasing in 1st variable.

2.36

Proposition :- Let  $I$  be fuzzy implication then

$$(i) \quad I(0, x) = 1 \quad \forall x \in [0, 1]$$

$$(ii) \quad I(x, 1) = 1 \quad \forall x \in [0, 1]$$

Proof. (i)

$$\text{Since } I(0, 0) = 1$$

$$1 = I(0, 0) \leq I(0, x)$$

$$\Rightarrow 1 \leq I(0, x) \quad \text{--- (i)}$$

$$\text{Also } I(0, x) \leq 1 \quad \text{--- (ii)}$$

$$\Rightarrow 1 \leq I(0, x) \leq 1$$

$$\Rightarrow I(0, x) = 1 \quad \forall x \in [0, 1]$$

$$(ii) \text{ Since } I(1, 1) = 1$$

$$1 = I(1, 1) \leq I(x, 1)$$

$$I(x, 1) \geq I(1, 1) = 1$$

$$I(x, 1) \geq 1 \quad \text{--- (i)}$$

Also

$$I(x, 1) \leq 1 \quad \text{--- (ii)}$$

$$1 \leq I(x, 1) \leq 1$$

$$\Rightarrow \boxed{I(x, 1) = 1} \quad \forall x \in [0, 1]$$

S-implication :-

Let  $S$  be t-conorm and  $N$  be a strong negation. Then

$$I(x, y) = S(N(x), y) \text{ is called}$$

S-implication.

Proposition 2.39 :-  $I(x, y) = S(N(x), y)$  is an <sup>above</sup> implication.

Proof 2-(i) Let  $x \leq z$  then  $N(x) \geq N(z)$

$$\text{and } S(N(x), y) \geq S(N(z), y)$$

$$I(x, y) \geq I(z, y)$$

So  $I$  is decreasing in 1st variable



$$y \leq z$$

$$S(N(x), y) \leq S(N(x), z)$$

$$= I(x, y) \leq I(x, z)$$

ph-33.

(ii) ~~Similarly~~ Similarly  $I$  is increasing in its 2nd variable.

$$(iii) \quad I(1, 0) = S(N(1), 0) = S(0, 0) = 0$$

$$I(1, 1) = S(N(1), 1) = S(0, 1) = 1$$

$$I(0, 0) = S(N(0), 0) = S(1, 0) = 1$$

Thus  $I(x, y) = S(N(x), y)$  is implication.

Example (2.4.1):-

$$\text{Let } S(x, y) = x + y - xy$$

$$N(x) = 1 - x$$

then  $I(x, y) = 1 - x + xy$  is the  $S$ -implication.

$S(x, y) = x + y - xy$  is  $t$ -conorm and  $N(x) = 1 - x$  is strong negation then  $S$ -implication is;

$$I(x, y) = S(N(x), y)$$

$$= S(1 - x, y)$$

$$= 1 - x + y - (1 - x)y$$

$$= 1 - x + y - y + xy$$

$$I(x, y) = 1 - x + xy \quad \text{is the } S\text{-implication.}$$

ion.

Example :- (2.4.2)

$$S(x, y) = \min(x + y, 1)$$

$$N(x) = 1 - x$$

$$I(x, y) = \min(1 - x + y, 1) \text{ is } S\text{-implication}$$

proof :-  $S$ -implication is;

$$I(x, y) = S(N(x), y)$$

$$= S(1 - x, y)$$

$$= \min(1 - x + y, 1) = I(x, y) \text{ proof.}$$

Residual-Implication (R-implication) :-

Let  $T$  be a t-norm.

$$I_T(x, y) = \text{Sup} \{ z \mid T(x, z) \leq y \}$$

is called R-implication.

Proposition (2.45) :-

R-implication is implication

Proof :-

$$(i) \quad x_1 \leq x_2 \Rightarrow T(x_1, z) \leq T(x_2, z), \forall z \in [0, 1]$$

$$\text{if } z_0 \in \{ z \mid T(x_2, z) \leq y \}$$

$$\Rightarrow z_0 \in \{ z \mid T(x_1, z) \leq y \}$$

$$\{ z \mid T(x_2, z) \leq y \} \subseteq \{ z \mid T(x_1, z) \leq y \}$$

$$I_T(x_2, y) = \text{Sup} \{ z \mid T(x_2, z) \leq y \} \leq \text{Sup} \{ z \mid T(x_1, z) \leq y \}$$

$$= I_T(x_1, y)$$

$$I_T(x_2, y) \leq I_T(x_1, y)$$

ii)

$$I_T(1, 0) = \text{Sup} \{ z \mid T(1, z) \leq 0 \}$$

$$= \text{Sup} \{ z \mid 0 \leq 0 \} =$$

$$= 0$$

$$I_T(1, 1) = \text{Sup} \{ z \mid T(1, z) \leq 1 \}$$

$$= \text{Sup} \{ z \mid 1 \leq 1 \} = \text{Sup} [0, 1]$$

$$= 1$$

$$I_T(0, 0) = \text{Sup} \{ z \mid T(0, z) \leq 0 \} =$$

$$=$$

(i)  
 Example 2.47 :-  $T(x, y) = \min(x, y)$  gives  
 the R-implication.

$$T(x, y) = x \wedge y \quad \text{then.}$$

$$I_T(x, y) = \text{Sup} \{z \mid x \wedge z \leq y\}$$

There are two cases if  $x \leq y$   
 then possible value of  $z$  is that at  
 which we take minimum with  $x$  and the  
 result is less or equal to  $y$ .

$$\rightarrow I_T(x, y) = \text{Sup} [0, 1] = 1.$$

if  $x > y$

$$I_T(x, y) = \text{Sup} \{z \mid x \wedge z \leq y\}$$

then possible value of  $z$  will be  
 less or equal to  $y$ .

$$I_T(x, y) = \text{Sup} [0, y] = y.$$

$$\text{Thus } I_T = \begin{cases} 1 & x \leq y \\ y & x > y \end{cases}$$

↑ example  $x = 0.2, y = 0.6$

if  $x < y$ ,

$$x \wedge z \leq 0.6$$

$$0.2 \wedge [0, 1] \leq 0.6$$

if  $x = 0.6, y = 0.2$

$x > y$

$$x \wedge z \leq y \Rightarrow 0.6 \wedge [0, 1] \leq 0.2$$

example 2.47 (ii)  $T(x, y) = xy$  gives the  $\mathcal{R}$ -implication.

$$T(x, y) = xy$$

$$I_T(x, y) = \text{Sup} \{z \mid xz \leq y\}$$

There are two cases

if  $x \leq y$

Then possible value of  $z \in [0, 1]$

$$I_T(x, y) = \text{Sup} \{[0, 1] \mid x[0, 1] \leq y\}$$

$$= \text{Sup}\{[0, 1]\} = 1$$

if  $x > y$

$$I_T(x, y) = \text{Sup} \{ [0, y/x] \mid xz \leq y \}$$

$$= \text{Sup} \{ [0, y/x] \mid xz \geq \frac{y}{x} \}$$

$$= \text{Sup} \{ [0, y/x] \} = y/x$$

$$I_T(x, y) = \begin{cases} 1 & \text{if } x \leq y \\ y/x & \text{if } x > y \end{cases}$$

example if  $x \leq y$

$$x = 0.2, y = 0.6$$

$$\rightarrow 0.2z \leq 0.6 \quad \text{if } z = 1$$

$$0.2 \leq 0.6 \quad \checkmark \quad z \in [0, 1]$$

if  $x > y$ ,  $x = 0.6, y = 0.2$

$$0.6 [0, \frac{0.2}{0.6}] \leq 0.2$$

$$0 \leq 0.2 \quad \checkmark \quad z = 0$$

$$0.2 < 0.2 \quad \checkmark \quad z = \frac{0.2}{0.6}$$

2.46  $T(x, 1) = x$

$$I_T = \text{Sup} \{z \mid T(x, 1) \leq y\}$$

$$= \text{Sup} \{z \mid x \leq y\} \Rightarrow I_T(x, y) = \text{Sup} \{[0, 1]\} = 1$$

### Fuzzy Equivalence :- (2.48)

A function  $E: [0, 1]^2 \rightarrow [0, 1]$  is a fuzzy equivalence if it satisfies the following conditions:

$$E_1: E(x, y) = E(y, x)$$

$$E_2: E(0, 1) = E(1, 0) = 0$$

$$E_3: E(x, x) = 1 \quad \forall x \in [0, 1]$$

$$E_4: \text{if } x \leq x' \leq y' \leq y \text{ then } E(x, y) \leq E(x', y')$$

denoted be  $x \leftrightarrow y$

Theorem 2.49:-

The following statements are equivalent:

(a)  $E$  is fuzzy equivalence

(b) There exist a fuzzy implication  $I$  with the property  $I(x, x) = 1 \quad \forall x \in [0, 1]$  such that

$$E(x, y) = \min [I(x, y), I(y, x)]$$

(b) There exist a fuzzy implication  $I$  with the property

$$E(x, y) = I(\max\{x, y\}, \min\{x, y\})$$

Proof :-  $b \rightarrow a$

Let  $E(x, y) = \min [I(x, y), I(y, x)]$ . Then we can easily check the condition properties  $E_1 \rightarrow E_3$ .

$$\begin{aligned} E_1: E(x, y) &= \min [I(x, y), I(y, x)] \\ &= \min [I(y, x), I(x, y)] \\ &= E(y, x) \end{aligned}$$

$$\Rightarrow E(x, y) = E(y, x) \quad \checkmark$$

$$\begin{aligned} E_2: E(0, 1) &= \min [I(0, 1), I(1, 0)] \\ &= \min [1, 0] = 0 \end{aligned}$$

$$E(1, 0) = \min [I(1, 0), I(0, 1)] = 0 \quad \checkmark$$

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$$\begin{aligned} E_3: E(x, x) &= \min [I(x, x), I(x, x)] \\ &= \min [1, 1] \quad \because I(x, x) = 1 \\ E(x, x) &= 1 \end{aligned}$$

$E_4$ : ~~Do~~ If  $x \leq x' \leq y' \leq y$  then  
 $E(x, y) \leq E(x', y')$

$$x \leq y$$

$$I(x, x) \geq I(y, x)$$

$$x \leq y$$

$$I(x, x) \leq I(x, y)$$

$$\Rightarrow I(y, x) \leq I(x, x) \leq I(x, y)$$

means  $E(x, y) = I(y, x)$

$$\Rightarrow E(x, y) = I(y, x)$$

Similarly  $E(x', y') = I(y', x')$

Now we show  $E(x, y) \leq E(x', y')$

or  $I(y, x) \leq I(y', x')$

$$x \leq x'$$

$$I(y', x) \leq I(y', x')$$

$$y \leq y'$$

$$I(y', x) \geq I(y, x)$$

$$\Rightarrow I(y, x) \leq I(y', x) \leq I(y', x')$$

$$\Rightarrow I(y, x) \leq I(y', x')$$

$$\Rightarrow E(x, y) \leq E(x', y')$$

Hence  $E$  is fuzzy equivalence.

Example  $I(x, y) = \min \{1-x+y, 1\}$

Find  $E(x, y)$  and then verify it is fuzzy equivalence

$$E(x, y) = \min [I(x, y), I(y, x)]$$

$$= \min [\min \{1-x+y, 1\}, \min \{1-y+x, 1\}]$$

$$E(x, y) = \min [1-x+y, 1-y+x]$$

(i)  $E(x, y) = E(y, x)$

$$E(x, y) = \min [1-x+y, 1-y+x]$$

$$= \min [1-y+x, 1-x+y]$$

$$E(x, y) = E(y, x) \quad \checkmark$$

(ii)  $E(0, 1) = E(1, 0) = 0$

$$E(0, 1) = \min [1-0+1, 1-1+0] = \min [2, 0]$$

$$E(0, 1) = 0$$

$$E(1, 0) = \min [1-1+0, 1-0+1] = 0 \quad \checkmark$$

(iii)  $E(x, x) = 1 \quad \forall x \in [0, 1]$

$$E(x, x) = \min [1-x+x, 1-x+x] = 1 \quad \checkmark$$

(iv) if  $x \leq x' \leq y' \leq y \Rightarrow E(x, y) \leq E(x', y')$

Since

$$I(x, y) = \min \{1-x+y, 1\}$$

if  $x \leq y$

$$\Rightarrow I(x, x) \geq I(y, x)$$

because  $1 \geq I(y, x)$

if  $x \leq y$

$$\Rightarrow I(x, x) \leq I(x, y)$$

$$\Rightarrow I(y, x) \leq I(x, x) \leq I(x, y)$$

$$\Rightarrow E(x, y) = I(y, x)$$

Similarly  $E(x', y') = I(y', x')$

Now we show  $I(y, x) \leq I(y', x')$

$$x \leq x'$$

$$I(y, x) \leq I(y', x')$$

$$y' \leq y$$

$$I(y', x) \geq I(y, x)$$

$$\Rightarrow I(y, x) \leq I(y', x) \leq I(y', x')$$

$$\Rightarrow I(y, x) \leq I(y', x')$$

$$\text{Thus } E(x, y) \leq E(x', y')$$

Hence  $E(x, y)$  is fuzzy Equivalence.

Example (Assignment) :-

$$I(x, y) = \begin{cases} 1 & \text{if } x \leq y \\ y/x & \text{if } x > y \end{cases}$$

Find Fuzzy equivalence and then verify.

$$\text{As } E(x, y) = \min [I(x, y), I(y, x)]$$

There are three cases

(1) if  $x = y$

$$E(x, y) = \min [1, 1] = 1$$

$$\text{as } I(y, x) = \begin{cases} 1 & \text{if } y \leq x \\ x/y & \text{if } y > x \end{cases}$$

(2) if  $x \leq y$

$$E(x, y) = \min [1, \overset{\min[I(x, y), I(y, x)]}{x/y}] = x/y$$

(3) if  $x > y$

$$E(x, y) = \min [y/x, 1] = y/x$$



$$\text{Thus } E(x, y) = \begin{cases} 1 & \text{if } x=y \\ x/y & \text{if } x < y \\ y/x & \text{if } y < x \end{cases}$$

Verification:-

$$(a) = E(x, y) = E(y, x)$$

$$E(x, y) = \begin{cases} 1 & \text{if } x=y \\ x/y & \text{if } x < y \\ y/x & \text{if } y < x \end{cases}$$

$$= \begin{cases} 1 & \text{if } y=x \\ y/x & \text{if } y < x \\ x/y & \text{if } x < y \end{cases}$$

$$E(x, y) = E(y, x)$$

$$(b) E(0, 1) = E(1, 0) = 0$$

$$E(0, 1) = \begin{cases} \frac{x}{y} & x < y \end{cases}$$

$$= \begin{cases} \frac{0}{1} & 0 < 1 \Rightarrow 0 \end{cases}$$

$$E(1, 0) = \begin{cases} \frac{y}{x} & \text{if } y < x \end{cases}$$

$$= \begin{cases} \frac{0}{1} & 0 < 1 \Rightarrow 0 \end{cases}$$

$$E(0, 1) = E(1, 0) = 0$$

$$(c) E(x, x) = 1 \quad \forall x \in [0, 1]$$

$$E(x, x) = \begin{cases} 1 & \text{if } y=x \\ x/x & \text{if } x < x \\ x/x & \text{if } x < x \end{cases}$$

$$= \begin{cases} 1 & \text{if } x=x = \boxed{1} \end{cases}$$

(d) if  $x \leq x' \leq y' \leq y \Rightarrow E(x, y) \leq E(x', y')$

$$E(x, y) = \min [I(x, y), I(y, x)]$$

$$x \leq y'$$

$$\Rightarrow I(x, x) \geq I(y, x)$$

$$\text{or } 1 \geq I(y, x)$$

$$x \leq y$$

$$\Rightarrow I(x, x) \leq I(x, y)$$

$$I(y, x) \leq I(x, x) \leq I(x, y)$$

$$\Rightarrow E(x, y) = I(y, x)$$

Also  $E(x', y') = I(y', x')$

Now we prove  $E(x, y) \leq E(x', y')$

$$x \leq x'$$

$$I(y', x) \leq I(y', x')$$

$$y' \leq y$$

$$I(y', x) \geq I(y, x)$$

$$I(y, x) \leq I(y', x) \leq I(y', x')$$

$$E(x, y) \leq E(x', y')$$

Thus  $E(x, y)$  is fuzzy equivalence

### Chapter 3 Fuzzy Relation.

classical Relation: A subset  $R \subset X \times Y$  where  $X$  and  $Y$  are classical sets is a classical relation.

It can be characterized by a function  $R: X \times Y \rightarrow \{0,1\}$

$$R(x,y) = \begin{cases} 1 & \text{if } (x,y) \in R \\ 0 & \text{otherwise} \end{cases}$$

### Fuzzy Relation:-

Let  $X$  and  $Y$  be two classical sets. A mapping  $R: X \times Y \rightarrow [0,1]$  is called a fuzzy relation. (OR)

A fuzzy relation on  $A \times B$  is denoted by  $R$  or  $R(x,y)$  is defined as the set

$$R(x,y) = \{((x,y), \mu_R(x,y)) \mid (x,y) \in A \times B, \mu_R(x,y) \in [0,1]\}$$

\*  $\mu_R(x,y) \in [0,1]$  is a degree of relationship between  $x$  and  $y$ .

\* We denote by  $F(X \times Y)$  the family of all fuzzy relations b/w elements of  $X$  and  $Y$ .

\* A fuzzy relation between elements in two finite sets  $X = \{x_1, x_2, \dots, x_m\}$  and  $Y = \{y_1, y_2, \dots, y_n\}$  can be represented as a matrix.

$$R = \begin{pmatrix} R(x_1, y_1) & R(x_1, y_2) & \dots & R(x_1, y_n) \\ R(x_2, y_1) & R(x_2, y_2) & \dots & R(x_2, y_n) \\ \vdots & \vdots & \ddots & \vdots \\ R(x_m, y_1) & R(x_m, y_2) & \dots & R(x_m, y_n) \end{pmatrix}$$

\* Since fuzzy relations are themselves fuzzy sets, it is possible to perform fuzzy sets operations on them

$$(i) N(R(x,y)) = 1 - R(x,y)$$

$$(ii) (R \vee S)(x,y) = R(x,y) \vee S(x,y)$$

$$(iii) (R \wedge S)(x,y) = R(x,y) \wedge S(x,y)$$

$$(iv) R^{-1}(x,y) = R(y,x)$$

We also consider t-norm and s-norm operations

$$(v) T(R, P)(x,y) = T(R(x,y), P(x,y))$$

$$(vi) S(R, P)(x,y) = S(R(x,y), P(x,y))$$

\* Example :- show how Fuzzy relation are represented  
 Let  $V = \{1, 2, 3\}$  and  $W = \{1, 2, 3\}$

A fuzzy relation  $R$  is a function defined in the space  $V \times W$  which takes values from the interval  $[0, 1]$ , expressed as  $R: V \times W \rightarrow [0, 1]$

$$R = [ \{ \{1, 1\}, 0.2 \}, \{ \{1, 2\}, 0.6 \}, \{ \{1, 3\}, 0.9 \}, \{ \{2, 1\}, 1 \}, \{ \{2, 2\}, 0.3 \}, \{ \{2, 3\}, 0.5 \}, \{ \{3, 1\}, 0 \}, \{ \{3, 2\}, 0.1 \}, \{ \{3, 3\}, 0.7 \} ]$$

In matrix form

$V \times W$	1	2	3
$R \triangleq 1$	0.2	0.6	0.9
2	1	0.3	0.5
3	0	0.1	0.7

\* **Composition** :- The operation "Composition" combine the fuzzy relation in different variables say  $(x, y)$  and  $(y, z)$ ;  
 $x \in A, y \in B, z \in C$ .

**Max-Min Composition** :- Let  $R \in F(X \times Y)$  and  $S \in F(Y \times Z)$  be fuzzy relations then  $R \circ S \in F(X \times Z)$ , defined as,

$$R \circ S(x, z) = \bigvee_{y \in Y} (R(x, y) \wedge S(y, z))$$
 is called the **Max-Min Composition** of  $R$  and  $S$ , denoted by  $R \circ S$ .

\* If  $R \in F(X \times Y)$  then we can define  $R^2 = R \circ R$  generally  $R^n = R \circ R^{n-1}$  for  $n \geq 2$

\* Let  $X = \{x_1, x_2, \dots, x_n\}$ ,  $Y = \{y_1, y_2, \dots, y_m\}$   
 $Z = \{z_1, z_2, \dots, z_p\}$  be finite sets.  
 if  $R = (r_{ij})_{i=1, \dots, n, j=1, \dots, m} \in F(X \times Y)$  and  $S = (s_{jk})_{j=1, \dots, m, k=1, \dots, p} \in F(Y \times Z)$  are fuzzy relations then the composition

$T = (t_{ik})_{i=1, \dots, n, k=1, \dots, p} \in F(X \times Z)$  is

$$t_{ik} = \bigvee_{j=1}^m (r_{ij} \wedge s_{jk})$$

**Example (3.7) page 35 :-**

if  $R = \begin{pmatrix} 0.3 & 0.7 & 0.2 \\ 1 & 0 & 0.9 \end{pmatrix}$

$S = \begin{pmatrix} 0.8 & 0.3 \\ 0.1 & 0 \\ 0.5 & 0.6 \end{pmatrix}$  then find  $R \circ S$

Solution: Consider

R $\triangleq$ X <sub>1</sub> Y	Y <sub>1</sub> Y <sub>2</sub> Y <sub>3</sub>			S $\triangleq$ Y <sub>1</sub> Y <sub>2</sub> Y <sub>3</sub>	Z <sub>1</sub> , Z <sub>2</sub>	
	Y <sub>1</sub>	Y <sub>2</sub>	Y <sub>3</sub>		Z <sub>1</sub>	Z <sub>2</sub>
X <sub>1</sub>	0.3	0.7	0.2	Y <sub>1</sub>	0.8	0.3
X <sub>2</sub>	1	0	0.9	Y <sub>2</sub>	0.1	0
				Y <sub>3</sub>	0.5	0.6

$$R \circ S(x, z) = V(R(x, y) \wedge S(y, z))$$

Step (1) Compute minimum operation

$$\min(R(x_1, y_1), S(y_1, z_1)) = \min(0.3, 0.8) = 0.3$$

$$\min(R(x_1, y_2), S(y_2, z_1)) = \min(0.7, 0.1) = 0.1$$

$$\min(R(x_1, y_3), S(y_3, z_1)) = \min(0.2, 0.5) = 0.2$$

$$\min(R(x_1, y_1), S(y_1, z_2)) = \min(0.3, 0.3) = 0.3$$

$$\min(R(x_1, y_2), S(y_2, z_2)) = \min(0.7, 0) = 0$$

$$\min(R(x_1, y_3), S(y_3, z_2)) = \min(0.2, 0.6) = 0.2$$

$$\min(R(x_2, y_1), S(y_1, z_1)) = \min(1, 0.8) = 0.8$$

$$\min(R(x_2, y_2), S(y_2, z_1)) = \min(0, 0.1) = 0$$

$$\min(R(x_2, y_3), S(y_3, z_1)) = \min(0.9, 0.5) = 0.5$$

$$\min(R(x_2, y_1), S(y_1, z_2)) = \min(1, 0.3) = 0.3$$

$$\min(R(x_2, y_2), S(y_2, z_2)) = \min(0, 0) = 0$$

$$\min(R(x_2, y_3), S(y_3, z_2)) = \min(0.9, 0.6) = 0.6$$

Step (2) Compute maximum operation:

$$\text{Max}(R(x_1, z_1)) = \max(0.3, 0.1, 0.2) = 0.3$$

$$\text{Max}(R(x_1, z_2)) = \max(0.3, 0, 0.2) = 0.3$$

$$\text{Max}(R(x_2, z_1)) = \max(0.8, 0, 0.5) = 0.8$$

$$\text{Max}(R(x_2, z_2)) = \max(0.3, 0, 0.6) = 0.6$$

Thus $R \circ S(x, z) = X \cdot Z$	$z_1$	$z_2$
$x_1$	0.3	0.3
$x_2$	0.8	0.6

is the required Max-Min Composition.

### Assignment No. 2

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proposition 3.8 :-

(i) The max-min Composition is associative

i.e  $(R \circ S) \circ Q = R \circ (S \circ Q)$

where  $R \in F(X \times Y)$ ,  $S \in F(Y \times Z)$ ,  $Q \in F(Z \times U)$ .

Proof :-

$$\begin{aligned} \text{LHS } (R \circ S) \circ Q(x, u) &= \max_z \{ \min \{ R \circ S(x, z), Q(z, u) \} \} \\ &= \max_z \{ \min \{ \max_y \{ \min \{ R(x, y), S(y, z) \} \}, Q(z, u) \} \} \end{aligned}$$

now for some  $y \in Y, k \in N$

$$= \max_z \{ \min \{ \min \{ R(x, y_k), S(y_k, z) \}, Q(z, u) \} \}$$

now for  $z_m \in Z, m \in N$

$$= \min \{ \min \{ R(x, y_k), S(y_k, z_m) \}, Q(z_m, u) \}$$

$$= \min \{ R(x, y_k), S(y_k, z_m), Q(z_m, u) \} \quad \text{--- (1)}$$

RHS

$$R \circ (S \circ Q)(x, u) = \max_y \{ \min \{ R(x, y), S \circ Q(y, u) \} \}$$

$$= \max_y \{ \min \{ R(x, y), \max_z \{ \min \{ S(y, z), Q(z, u) \} \} \} \}$$

$$= \max_y \{ \min \{ R(x, y), \min \{ S(y, z_m), Q(z_m, u) \} \} \}$$

for  $z_m \in Z, m \in N$

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$$= \min \{ R(x, y_k), \min \{ S(y_k, z_m), Q(z_m, u) \} \}$$

for  $y_k \in Y, k \in \mathbb{N}$

$$= \min \{ R(x, y_k), S(y_k, z_m), Q(z_m, u) \} \quad \text{--- (2)}$$

From (1) and (2) we have

$$\text{LHS} = \text{RHS}$$

$$(R \circ S) \circ Q = R \circ (S \circ Q)$$

(ii) Let  $R_1, R_2 \in F(X \times Y)$  and  $Q \in F(Y \times Z)$ .

If  $R_1 \leq R_2$  then  $R_1 \circ Q \leq R_2 \circ Q$

Proof :-

$$R_1 \circ Q(x, z) = \bigvee_y \{ R_1(x, y) \wedge Q(y, z) \}$$

$$R_2 \circ Q(x, z) = \bigvee_y \{ R_2(x, y) \wedge Q(y, z) \}$$

$$\bigvee_y \{ R_1(x, y) \wedge Q(y, z) \} \leq \bigvee_y \{ R_2(x, y) \wedge Q(y, z) \}$$

$$\Rightarrow R_1 \circ Q(x, z) \leq R_2 \circ Q(x, z)$$

Proposition 3.9 :- For any  $R, S \in F(X \times Y)$  and  $Q \in F(Y \times Z)$  we have

$$(i) (R \vee S) \circ Q = (R \circ Q) \vee (S \circ Q)$$

proof :-

$$(R \vee S) \circ Q = \bigvee_{y \in Y} \{ (R \vee S)(x, y) \wedge Q(y, z) \}$$

$$= \bigvee_{y \in Y} (R(x, y) \wedge Q(y, z)) \vee (S(x, y) \wedge Q(y, z))$$

$$\leq \bigvee_{y \in Y} R(x, y) \wedge Q(y, z) \vee \bigvee_{y \in Y} S(x, y) \wedge Q(y, z)$$



$$(R \vee S) \circ Q \leq (R \circ Q) \vee (S \circ Q) \text{ --- (1)}$$

or the other have  $R \circ Q \leq (R \vee S) \circ Q$

and  $S \circ Q \leq (R \vee S) \circ Q$

$$\Rightarrow (R \circ Q) \vee (S \circ Q) \leq (R \vee S) \circ Q \text{ --- (2)}$$

From (1) and (2) we have

$$(R \vee S) \circ Q = (R \circ Q) \vee (S \circ Q)$$

$$(ii) (R \wedge S) \circ Q \leq (R \circ Q) \wedge (S \circ Q)$$

$$(R \wedge S) \circ Q(x, z) = \bigvee_{y \in Y} (R \wedge S)(x, y) \wedge Q(y, z)$$

$$= \bigvee_{y \in Y} (R(x, y) \wedge S(x, y)) \wedge (S(x, y) \wedge Q(y, z))$$

$$\leq \bigvee_{y \in Y} (R(x, y) \wedge Q(y, z)) \wedge \bigvee_{y \in Y} (S(x, y) \wedge Q(y, z))$$

$$(R \wedge S) \circ Q(x, z) \leq (R \circ Q) \wedge (S \circ Q)$$

Min-Max Composition: Let  $R \in F(X \times Y)$  and  $S \in F(Y \times Z)$ . then  $R \circ S \in F(X \times Z)$ . defined as

$R \circ S(x, z) = \bigwedge_{y \in Y} R(x, y) \vee S(y, z)$  is called min-max composition of the fuzzy relation  $R$  and  $S$ .

Example (3.12) if  $R = \begin{pmatrix} 0.3 & 0.7 & 0.2 \\ 1 & 0 & 0.9 \end{pmatrix}$ ,  $S = \begin{pmatrix} 0.8 & 0.3 \\ 0.1 & 0 \\ 0.5 & 0.6 \end{pmatrix}$

Find  $R \circ S$ .

Solution:

(i) Compute max operation:

$$\max [R(x_1, y_1), S(y_1, z_1)] = \max [0.3, 0.8] = 0.8$$

$$\max [R(x_1, y_2), S(y_2, z_1)] = \max [0.7, 0.1] = 0.7$$

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$$\max [R(x_1, y_3), S(y_3, z_1)] = \max [0.2, 0.5] = 0.5$$

$$= \max [0.3, 0.3] = 0.3$$

$$= \max [0.7, 0] = 0.7$$

$$= \max [0.2, 0.6] = 0.6$$

$$= \max [1, 0.8] = 1$$

$$= \max [0, 0.1] = 0.1$$

$$= \max [0.9, 0.5] = 0.9$$

$$= \max [1, 0.3] = 1$$

$$= \max [0, 0] = 0$$

$$= \max [0.9, 0.6] = 0.9$$

(ii) Compute minimum minimum operation:-

$$R \circ S(x, z) = \min [R(x, y_1)] = \min [0.8, 0.7, 0.5]$$

$$= 0.5$$

$$= \min [0.3, 0.7, 0.6] = 0.3$$

$$= \min [1, 0.1, 0.9] = 0.1$$

$$= \min [1, 0, 0.9] = 0$$

$R \circ S(x, z) = x \times z$	$z_1$	$z_2$
$x_1$	0.5	0.3
$x_2$	0.1	0

Assignment (02) Umar Saeed Roll no. 29.

Proposition (3.13): (i) The min-max Composition is associative i.e.  $(R \circ S) \circ T = R \circ (S \circ T)$   
 For  $R \in F(X, Y)$ ,  $S \in F(Y, Z)$ ,  $T \in F(Z, U)$

Proof :- LHS

$$(R \circ S) \circ T = \min_z \{ \max \{ R \circ S(x, z), T(z, u) \} \}$$

$$= \min_z \{ \max \{ \min_y \{ \max \{ R(x, y), S(y, z) \}, T(z, u) \} \}$$

$$= \min_z \{ \max \{ \max \{ R(x, y_k), S(y_k, z) \}, T(z, u) \} \}$$

for some  $y_k \in Y$  and  $k \in \mathbb{N}$

$$= \max \{ \max \{ R(x, y_k), S(y_k, z_m) \}, T(z_m, u) \}$$

for some  $z_m \in Z$  and  $m \in \mathbb{N}$

$$= \max \{ R(x, y_k), S(y_k, z_m), T(z_m, u) \} \quad \text{--- (1)}$$

RHS

$$R \circ (S \circ T)(x, u) = \min_y \{ \max \{ R(x, y), S \circ T(y, u) \} \}$$

$$= \min_y \{ \max \{ R(x, y), \min_z \{ \max \{ S(y, z), T(z, u) \} \} \}$$

$$= \min_y \{ \max \{ R(x, y), \max \{ S(y, z_m), T(z_m, u) \} \} \}$$

for some  $z_m \in Z$ ,  $m \in \mathbb{N}$

$$= \max \{ R(x, y_k), \max \{ S(y_k, z_m), T(z_m, u) \} \}$$

for some  $y_k \in Y$ ,  $k \in \mathbb{N}$

$$= \max \{ R(x, y_k), S(y_k, z_m), T(z_m, u) \} \quad \text{--- (2)}$$

From (1) and (2) we have

$$(R \circ S) \circ T = R \circ (S \circ T)$$

(ii) Consider  $R_1, R_2 \in F(X \times Y)$  and  $Q \in F(Y \times Z)$ .

if  $R_1 \leq R_2$  then

$$R_1 \circ Q \leq R_2 \circ Q$$

Proof :-

$$R_1 \circ Q = \bigwedge_{y \in Y} R_1(x, y) \vee Q(y, z)$$

$$\leq \bigwedge_{y \in Y} R_2(x, y) \vee Q(y, z) = R_2 \circ Q$$

$$\Rightarrow R_1 \circ Q(x, z) \leq R_2 \circ Q(x, z) \text{ if } R_1 \leq R_2$$

Proposition 3.14 :- For any  $R, S \in F(X \times Y)$

and  $T \in F(Y \times Z)$  we have

$$(i) (R \wedge S) \circ T = (R \circ T) \wedge (S \circ T)$$

$$(ii) (R \vee S) \circ T \geq (R \circ T) \vee (S \circ T)$$

Proof :-

$$(i) (R \wedge S) \circ T = \bigwedge_{y \in Y} ((R \wedge S) \vee T(y, z))$$

$$= \bigwedge_{y \in Y} [R(x, y) \vee T(y, z) \wedge S(x, y) \vee T(y, z)]$$

$$\geq [\bigwedge_{y \in Y} R(x, y) \vee T(y, z)] \wedge [\bigwedge_{y \in Y} S(x, y) \vee T(y, z)]$$

$$\geq R \circ T \wedge S \circ T$$

$$\geq (R \circ T) \wedge (S \circ T) \text{ --- (1)}$$

on the other hand

$$R \circ T \geq (R \wedge S) \circ T$$

$$S \circ T \geq (R \wedge S) \circ T$$

$$\Rightarrow (R \circ T) \wedge (S \circ T) \geq (R \wedge S) \circ T \text{ --- (2)}$$

From (1) and (2)

$$(R \wedge S) \circ T = (R \circ T) \wedge (S \circ T)$$

$$\begin{aligned}
 \text{(ii)} \quad (R \vee S) \circ T &= \bigwedge_{y \in Y} (R \vee S) \vee T(y, z) \\
 &= \bigwedge_{y \in Y} (R(x, y) \vee T(y, z)) \vee (S(x, y) \vee T(y, z)) \\
 &\geq \bigwedge_{y \in Y} [R(x, y) \vee T(y, z)] \vee \bigwedge_{y \in Y} [S(x, y) \vee T(y, z)] \\
 &= (R \circ T) \vee (S \circ T)
 \end{aligned}$$

$$(R \vee S) \circ T \geq (R \circ T) \vee (S \circ T)$$

Remark (3.10) :-  $R = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$ ,  $S = \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}$

$$Q = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

$$\Rightarrow (R \wedge S) \circ Q = \begin{pmatrix} 0 & 0 \\ 1 & 1 \end{pmatrix}$$

and  $(R \circ Q) \wedge (S \circ Q) = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$

Remark (3.15) :-  $R = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$ ,  $S = \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}$

$$T = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}$$

$$\Rightarrow (R \vee S) \circ T = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

and  $(R \circ T) \vee (S \circ T) = \begin{pmatrix} 0 & 0 \\ 1 & 1 \end{pmatrix}$

Proposition (3.16) :- If we consider the standard negative we have (i)  $\overline{R \circ S} = \overline{R} \circ \overline{S}$ , (ii)  $\overline{R \circ S} = \overline{R} \circ \overline{S}$

(i)  $\overline{R \circ S}(x, z) = \bigvee_{y \in Y} \overline{R(x, y) \wedge S(y, z)}$

$$= \bigwedge_{y \in Y} \overline{R(x, y) \wedge S(y, z)}$$

$$= \bigwedge_{y \in Y} \overline{R(x, y)} \vee \overline{S(y, z)} = \overline{R} \circ \overline{S}$$

(ii) similar way.

### Min $\rightarrow$ Composition:

Let  $\rightarrow$  be the standard Gödel implication defined as

$$x \rightarrow y = \sup \{z \in [0,1] \mid x \wedge z \leq y\} = \begin{cases} 1 & \text{if } x \leq y \\ y & \text{otherwise} \end{cases}$$

Proposition 3.18: For any  $x, y, z \in [0,1]$  we have

(i)  $(x \vee y) \rightarrow z = (x \rightarrow z) \wedge (y \rightarrow z)$

Proof:

Case (1): if  $x \vee y \leq z$  then  $x \vee y \rightarrow z = 1$  — (A)

also we have  $x \leq z$  and  $y \leq z$

then  $x \rightarrow z = 1, y \rightarrow z = 1$

$$(x \rightarrow z) \wedge (y \rightarrow z) = 1 \wedge 1 = 1 \text{ — (B)}$$

For  $A=B$

Case (2): if  $x \vee y > z$  then  $(x \vee y) \rightarrow z = z$  — (C)

also we have  $x \rightarrow z = z, y \rightarrow z = z$

$$\text{Thus } (x \rightarrow z) \wedge (y \rightarrow z) = z \text{ — (D)}$$

Thus in both cases

$$(x \vee y) \rightarrow z = (x \rightarrow z) \wedge (y \rightarrow z)$$

(ii)  $(x \wedge y) \rightarrow z = (x \rightarrow z) \vee (y \rightarrow z)$

Case (1): if  $x \wedge y \leq z$  then

$$(x \wedge y) \rightarrow z = 1 \text{ — (A) Also } x \leq z \text{ and } y \leq z$$

then  $x \rightarrow z = 1, y \rightarrow z = 1$

$$\text{so } (x \rightarrow z) \vee (y \rightarrow z) = 1 \text{ — (B)}$$

$A=B$

Case (2): if  $x \wedge y > z$  then

$$x \wedge y \rightarrow z = z \text{ — (C) also } x > z, y > z$$

then  $x \rightarrow z = z, y \rightarrow z = z$

$$\text{so } (x \rightarrow z) \vee (y \rightarrow z) = z \text{ — (D)}$$

Thus in both cases

$$(x \wedge y) \rightarrow z = (x \rightarrow z) \vee (y \rightarrow z)$$

$$(iii) \quad x \rightarrow (y \vee z) = (x \rightarrow y) \vee (x \rightarrow z)$$

Case (1) if  $x \leq y \vee z$  then  $x \rightarrow y \vee z = 1$

also  $x \leq y$  <sup>or</sup> ~~and~~  $x \leq z$  then

$$(x \rightarrow y) \vee (x \rightarrow z) = 1 \vee 1 = 1$$

Case (2): if  $x > y \vee z$  then

$$x \rightarrow y \vee z = y \vee z$$

also  $x > y \Rightarrow x \rightarrow y = y$

or  $x > z \Rightarrow x \rightarrow z = z$

$$(x \rightarrow y) \vee (x \rightarrow z) = y \vee z$$

Thus in Both cases

$$x \rightarrow (y \vee z) = (x \rightarrow y) \vee (x \rightarrow z)$$

$$(iv) \quad x \rightarrow (y \wedge z) = (x \rightarrow y) \wedge (x \rightarrow z)$$

Case (1) if  $x \leq y \wedge z$  then

$$x \rightarrow (y \wedge z) = 1 \quad \text{also}$$

$x \leq y$  or  $x \leq z$  then  $x \rightarrow y = 1$  or  $x \rightarrow z = 1$

$$(x \rightarrow y) \wedge (x \rightarrow z) = 1 \wedge 1 = 1$$

Case (2) if  $x > y \wedge z$  then

$$x \rightarrow y \wedge z = y \wedge z \quad \text{also}$$

$x > y \Rightarrow x \rightarrow y = y$  and or

$x > z \Rightarrow x \rightarrow z = z$  thus

$$(x \rightarrow y) \wedge (x \rightarrow z) = y \wedge z$$

In Both cases

$$x \rightarrow (y \wedge z) = (x \rightarrow y) \wedge (x \rightarrow z)$$

$$(v) \quad x \wedge (x \rightarrow y) \leq y$$

Case (1) if  $x \leq y$  then  $x \wedge (x \rightarrow y) = x \wedge 1$   
 $= x \leq y$  (already)

Case (2) if  $x > y$  then  $x \wedge (x \rightarrow y)$

$$= x \wedge y = y$$

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In: Both cases

$$x \wedge (x \rightarrow y) \leq y$$

$$(iv) \quad x \rightarrow (x \wedge y) \geq y$$

Case (1) if  $x \leq y$  then  $x \rightarrow x \wedge y = x \rightarrow x = 1 \geq y$

Case (2) if  $x > y$

$$x \rightarrow (x \wedge y) = \cancel{x \wedge y} \quad x \rightarrow y = y$$

Thus  $x \rightarrow (x \wedge y) \geq y$

(vii)

$$(x \rightarrow y) \rightarrow y \geq x$$

Case (1) if  $x \leq y$

$$(x \rightarrow y) \rightarrow y = 1 \rightarrow y = y \geq x \quad (\text{already})$$

because  $1 \neq y$

Case (2)  $x > y$

$$(x \rightarrow y) \rightarrow y = y \rightarrow y = 1 \geq x$$

$$\Rightarrow (x \rightarrow y) \rightarrow y \geq x$$



In the present <sup>work</sup> we <sup>mainly</sup> ~~simply~~ use the subcomposition and we call it simply  $\rightarrow$  min  $\rightarrow$  Composition.

Proposition 3.21:- If  $R, S \in F(X \times Y)$  and  $Q \in F(Y \times Z)$  are such that  $R \leq S$  then

$$R \triangleleft Q(n, z) \geq S \triangleleft Q(n, z)$$

Since  $\rightarrow$  is decreasing in the first argument we have

$$R \triangleleft Q(n, z) = \bigwedge_{y \in Y} R(n, y) \rightarrow Q(y, z)$$

$$\geq \bigwedge_{y \in Y} S(n, y) \rightarrow Q(y, z) = S \triangleleft Q(n, z)$$

Proposition 3.22:- For any  $R, S \in F(X \times Y)$  and  $Q \in F(Y \times Z)$  we have

$$(i) (R \vee S) \triangleleft Q = (R \triangleleft Q) \wedge (S \triangleleft Q)$$

From previous proposition

$$(R \vee S) \triangleleft Q \leq R \triangleleft Q \quad \text{and}$$

$$(R \vee S) \triangleleft Q \leq S \triangleleft Q$$

thus

$$(R \vee S) \triangleleft Q \leq (R \triangleleft Q) \wedge (S \triangleleft Q) \quad (a)$$

Also

$$(R \vee S) \triangleleft Q(n, z) = \bigwedge_{y \in Y} ((R(n, y) \vee S(n, y)) \rightarrow Q(y, z))$$

$$= \bigwedge_{y \in Y} (R(n, y) \rightarrow Q(y, z)) \wedge (S(n, y) \rightarrow Q(y, z))$$

$$\geq \bigwedge_{y \in Y} (R(n, y) \rightarrow Q(y, z)) \wedge \bigwedge_{y \in Y} (S(n, y) \rightarrow Q(y, z))$$

$$= (R \triangleleft Q) \wedge (S \triangleleft Q) \quad (b)$$

from (a) and (b).

$$(R \vee S) \triangleleft Q = (R \triangleleft Q) \wedge (S \triangleleft Q)$$

$$(ii) (R \wedge S) \triangleleft Q \geq (R \triangleleft Q) \vee (S \triangleleft Q)$$

as

$$(R \wedge S) \triangleleft Q \geq (R \triangleleft Q)$$

and

$$(R \wedge S) \triangleleft Q \geq (S \triangleleft Q)$$

thus

$$(R \wedge S) \triangleleft Q \geq (R \triangleleft Q) \vee (S \triangleleft Q)$$

Fuzzy Relational Eq's with Max-Min and

Min  $\rightarrow$  Compositions :- We consider the following two fuzzy relational eq,

$$R \circ P = Q$$

and

$$R \triangleleft P = Q$$

with

$$R \in F(X \times Y), P \in F(Y \times Z) \quad \text{and}$$

$$Q \in F(X \times Z)$$

P.T.O

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Assignment : No. 02

Umar Saeed Roll no. 29

Theorem 3.23 :- The following inequalities hold true.

(i)  $P \leq R' \triangleleft (R \circ P)$

Proof :- For every  $y \in Y$  and  $z \in Z$  we have

$$R' \triangleleft (R \circ P)(y, z) = \bigwedge_{x \in X} R'(y, x) \rightarrow (R \circ P)(x, z)$$

$$= \bigwedge_{x \in X} R'(y, x) \rightarrow \bigvee_{t \in Y} R(x, t) \wedge P(t, z)$$

$$\geq \bigwedge_{x \in X} R(x, y) \rightarrow R(x, y) \wedge P(y, z)$$

$$\geq \bigwedge_{x \in X} P(y, z) = P(y, z)$$

$$\therefore x \rightarrow y = 1 \text{ if } x \leq y$$

$$x \rightarrow y = y \text{ if } x > y$$

$$\Rightarrow P \leq R' \triangleleft (R \circ P)$$

(ii)  $R \circ (R' \triangleleft Q) \leq Q$

Proof :- for every  $x \in X$  and  $z \in Z$  we have

$$R \circ (R' \triangleleft Q)(x, z) = \bigvee_{y \in Y} R(x, y) \wedge (R' \triangleleft Q)(y, z)$$

$$= \bigvee_{y \in Y} R(x, y) \wedge \left( \bigwedge_{t \in X} R'(y, t) \rightarrow Q(t, z) \right)$$

$$\leq \bigvee_{y \in Y} R(x, y) \wedge (R'(y, x) \rightarrow Q(x, z))$$

$$\leq \bigvee_{y \in Y} R(x, y) \wedge (R(x, y) \rightarrow Q(x, z))$$

$$\leq \bigvee_{y \in Y} R(x, y) \wedge Q(x, z) \leq \bigvee_{y \in Y} Q(x, z) = Q(x, z)$$

$$\therefore x \rightarrow y = y \text{ if } x > y$$

$$\therefore x \rightarrow y = 1 \text{ if } x \leq y$$

Thus

$$R \circ (R' \triangleleft Q)(x, z) \leq Q(x, z)$$

$$(iii) R \leq (P \triangleleft (R \circ P)^{-1})^{-1}$$

Proof: For every  $x \in X, y \in Y$

$$(P \triangleleft (R \circ P)^{-1})^{-1}(x, y) = P \triangleleft (R \circ P)^{-1}(y, x)$$

$$= \bigwedge_{z \in Z} P(y, z) \rightarrow (R \circ P)^{-1}(z, x)$$

$$= \bigwedge_{z \in Z} P(y, z) \rightarrow (R \circ P)(x, z)$$

$$= \bigwedge_{z \in Z} P(y, z) \rightarrow \bigvee_{t \in Y} R(x, t) \wedge P(t, z)$$

$$\geq \bigwedge_{z \in Z} P(y, z) \rightarrow R(x, y) \wedge P(y, z)$$

$$\geq \bigwedge_{z \in Z} R(x, y) \wedge P(y, z) \geq \bigwedge_{z \in Z} R(x, y) = R(x, y)$$

$$(iv) (P \triangleleft Q^{-1})^{-1} \circ P \leq Q$$

Proof: For every  $x \in X, z \in Z$  we have

$$(P \triangleleft Q^{-1})^{-1} \circ P(x, z) = \bigvee_{y \in Y} (P \triangleleft Q^{-1})^{-1}(x, y) \wedge P(y, z)$$

$$= \bigvee_{y \in Y} (P \triangleleft Q^{-1})(y, x) \wedge P(y, z)$$

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$$= \bigvee_{y \in Y} \left( \bigwedge_{t \in Z} P(y, t) \rightarrow Q'(t, y) \right) \wedge P(y, z)$$

$$\leq \bigvee_{y \in Y} \left( \bigwedge_{z \in Z} P(y, z) \rightarrow Q'(z, y) \right) \wedge P(y, z)$$

$$\leq \bigvee_{y \in Y} \left( \bigwedge P(y, z) \rightarrow Q(x, z) \right) \wedge P(y, z)$$

$$\leq \bigvee_{y \in Y} Q(x, z) \wedge P(y, z) \leq \bigvee_{y \in Y} Q(x, z)$$

Thus

$$(P \triangleleft Q') \circ P \leq Q$$

Proposition: 3.26 :- Following inequality holds true:

$$(i) (Q \triangleleft P') \triangleleft P \geq Q$$

Proof :- taking  $\rightarrow$  decreasing in its first arguments we have

$$(Q \triangleleft P') \triangleleft P(x, z) = \bigwedge_{y \in Y} \left( (Q \triangleleft P')(y, y) \rightarrow P(y, z) \right)$$

$$= \bigwedge_{y \in Y} \left( \bigwedge_{t \in Z} Q(x, t) \rightarrow P'(x, t) \right) \rightarrow P(y, z)$$

$$= \bigwedge_{y \in Y} \left( \bigwedge_{t \in Z} Q(x, t) \rightarrow P(t, y) \right) \rightarrow P(y, z)$$

$$\geq \bigwedge_{y \in Y} \left( Q(x, z) \rightarrow P(y, z) \right) \rightarrow P(y, z)$$

$$\geq \bigwedge_{y \in Y} \left( Q(x, z) = Q(x, z) \right)$$

Thus  $(Q \triangleleft P') \triangleleft P(x, z) \geq Q$

$$(ii) (R \triangleleft P) \triangleleft \bar{P}' \geq R$$

Proof :-

$$\begin{aligned} (R \triangleleft P) \triangleleft \bar{P}'(x,y) &= \bigwedge_{z \in Z} (R \triangleleft P) \xrightarrow{(x,z)} \bar{P}'(z,y) \\ &= \bigwedge_{z \in Z} (R \triangleleft P)(x,z) \rightarrow \bar{P}'(z,y) \\ &= \bigwedge_{z \in Z} \left( \bigwedge_{t \in Y} R(x,t) \rightarrow P(t,z) \right) \rightarrow \bar{P}'(z,y) \\ &\geq \bigwedge_{z \in Z} (R(x,y) \rightarrow P(y,z)) \rightarrow P(y,z) \\ &\geq \bigwedge_{z \in Z} R(x,y) \end{aligned}$$

Thus  $(R \triangleleft P) \triangleleft \bar{P}' \geq R$

$$(iii) R^{-1} \circ (R \triangleleft P) \leq P$$

Proof :-

$$\begin{aligned} R^{-1} \circ (R \triangleleft P)(y,z) &= \bigvee_{x \in X} R^{-1}(y,x) \wedge (R \triangleleft P)(x,z) \\ &= \bigvee_{x \in X} R(x,y) \wedge \left( \bigwedge_{t \in Y} R(x,t) \rightarrow P(t,z) \right) \\ &\leq \bigvee_{x \in X} R(x,y) \wedge (R(x,y) \rightarrow P(y,z)) \end{aligned}$$

From proposition 3.18

$$x \wedge (x \rightarrow y) \leq y$$

$$\leq \bigvee_{x \in X} P(y,z) = P(y,z)$$

Thus  $R^{-1} \circ (R \triangleleft P)(y,z) \leq P(y,z)$

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$$(iv) R \triangleleft (R^{-1} \circ Q) \geq Q$$

Proof:

$$R \triangleleft (R^{-1} \circ Q)(x, z) = \bigwedge_{y \in Y} R(x, y) \rightarrow (R^{-1} \circ Q)(y, z)$$

$$= \bigwedge_{y \in Y} R(x, y) \rightarrow \left( \bigvee_{t \in X} (R^{-1}(y, t) \wedge Q(t, z)) \right)$$

$$\geq \bigwedge_{y \in Y} R(x, y) \rightarrow (R(x, y) \wedge Q(x, z))$$

From proposition 3.18

$$x \rightarrow (x \wedge y) \geq y$$

$$\geq \bigwedge_{y \in Y} Q(x, z) = Q(x, z)$$

Thus

$$R \triangleleft (R^{-1} \circ Q) \geq Q$$

Theorem 3.27:

(i) Consider the equation with  $R \triangleleft P = Q$  with unknown  $R$ . The equation has a solution iff  $Q \triangleleft P^{-1}$  is a solution, and in this case it is the greatest solution of this Eq.

(ii) Consider the eq.  $R \triangleleft P = Q$  with unknown  $P$ . The equation has a solution iff  $R^{-1} \circ Q$  is a solution, and in this case it is the least solution of this eq.